Experimental Mathematics & Computer Algebra

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Part II

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Short Summary of Part I

Experimental Mathematics

A 3-step process:

- 1. Compute a high order approximation (high precision numerical approx.,
 - power series truncated to high order, large number of terms in a sequence,...)
- 2. Guess/conjecture a general formula (with the help of a computer)
- 3. Prove it *(using computer-algebra algorithms)*

Basic Computer Algebra

Simplification is undecidable.

Fast computation with large integers, polynomials, power series, matrices,...

Tools for conjectures:

OEIS (integer sequences)
ISC (real numbers)

Padé-Hermite approximants (relations between power series) LLL (relations between numerical approximations).

Exercises for Part I

Ex 1. Identities for Nice Constants

$$\sum_{i=0}^{\infty} \frac{16^{-i}}{8i+j}, \qquad j = 1, \dots, 8$$

lead to formulas for

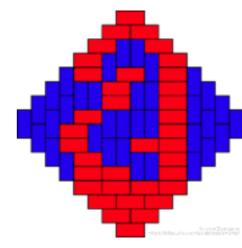
 π , ln 2, ln 3, ln 5, arctan 2, arctan 3, $\sqrt{2}$ arctan $(1/\sqrt{2})$, $\sqrt{2}$ ln $(1 + \sqrt{2})$.

Ex 2. Stanley's Problem E 2297

Number of monomials in a generic nxn symmetric matrix

Ex 3. Number of Domino Tilings of the Aztec Diamond

Guess & Prove a formula



Part II: Tools for Proofs

A short tour of ``univariate'' computer algebra

- 1. Resultants
- 2. D-finite sequences & series
- 3. Creative telescoping

1. Resultants

Polynomials as a data-structure

Definition

The Sylvester matrix of $A = a_m x^m + \cdots + a_0 \in \mathbb{K}[x]$, $(a_m \neq 0)$, and of $B = b_n x^n + \cdots + b_0 \in \mathbb{K}[x]$, $(b_n \neq 0)$, is the square matrix of size m + n

The resultant Res(A, B) of A and B is the determinant of Syl(A, B).

▶ Definition extends to polynomials with coefficients in a commutative ring R.

Basic Observation

If
$$A = a_m x^m + \dots + a_0$$
 and $B = b_n x^n + \dots + b_0$, then

$$\begin{bmatrix} a_m & a_{m-1} & \dots & a_0 \\ & \ddots & \ddots & & \ddots \\ & & a_m & a_{m-1} & \dots & a_0 \\ b_n & b_{n-1} & \dots & b_0 \\ & & \ddots & \ddots & & \ddots \\ & & b_n & b_{n-1} & \dots & b_0 \end{bmatrix} \times \begin{bmatrix} \alpha^{m+n-1} \\ \vdots \\ \alpha \\ 1 \end{bmatrix} = \begin{bmatrix} \alpha^{n-1}A(\alpha) \\ \vdots \\ A(\alpha) \\ \alpha^{m-1}B(\alpha) \\ \vdots \\ B(\alpha) \end{bmatrix}$$

Corollary: If $A(\alpha) = B(\alpha) = 0$, then Res(A, B) = 0.

Example: the Discriminant

The discriminant of A is the resultant of A and of its derivative A'.

E.g. for
$$A = ax^{2} + bx + c$$
,

$$\mathsf{Disc}(A) = \mathsf{Res}\,(A,A') = \det \left[\begin{array}{ccc} a & b & c \\ 2a & b \\ & 2a & b \end{array} \right] = -a(b^2 - 4ac).$$

E.g. for
$$A = ax^3 + bx + c$$
,

The discriminant vanishes when A and A' have a common root, that is when A has a multiple root.

Main Properties

• Link with gcd $\operatorname{Res}(A, B) = 0$ if and only if $\gcd(A, B)$ is non-constant.

Elimination property

There exist $U, V \in \mathbb{K}[x]$ not both zero, with $\deg(U) < n$, $\deg(V) < m$ and such that the following Bézout identity holds:

$$Res(A, B) = UA + VB$$
 in $\mathbb{K} \cap (A, B)$.

Poisson formula

If
$$A = a(x - \alpha_1) \cdots (x - \alpha_m)$$
 and $B = b(x - \beta_1) \cdots (x - \beta_n)$, then
$$\operatorname{Res}(A, B) = a^n b^m \prod_{i,j} (\alpha_i - \beta_j) = a^n \prod_{1 \le i \le m} B(\alpha_i).$$

Bézout-Hadamard bound

If $A, B \in \mathbb{K}[x, y]$, then $\operatorname{Res}_y(A, B)$ is a polynomial in $\mathbb{K}[x]$ of degree $\leq \deg_x(A) \deg_y(B) + \deg_x(B) \deg_y(A)$.

Application: Computation with Algebraic Numbers

Let
$$A=\prod_i(x-\alpha_i)$$
 and $B=\prod_j(x-\beta_j)$ be polynomials of $\mathbb{K}[x]$. Then
$$\operatorname{Res}_x(A(x),B(t-x))=\prod_{i,j}(t-(\alpha_i+\beta_j)),$$

$$\operatorname{Res}_x(A(x),B(t+x))=\prod_{i,j}(t-(\beta_j-\alpha_i)),$$

$$\operatorname{Res}_x(A(x),x^{\deg B}B(t/x))=\prod_{i,j}(t-\alpha_i\beta_j),$$

$$\operatorname{Res}_x(A(x),t-B(x))=\prod_i(t-B(\alpha_i)).$$

In particular, the set of algebraic numbers is a field.

Proof: Poisson's formula. E.g., first one:
$$\prod_{i} B(t - \alpha_i) = \prod_{i,j} (t - \alpha_i - \beta_j).$$

► The same formulas apply mutatis mutandis to algebraic power series.

Exercise for the afternoon: A Nice Number

Guess and prove a simple formula for

$$\frac{\sin\frac{2\pi}{7}}{\sin^2\frac{3\pi}{7}} - \frac{\sin\frac{\pi}{7}}{\sin^2\frac{2\pi}{7}} + \frac{\sin\frac{3\pi}{7}}{\sin^2\frac{\pi}{7}}.$$

Rothstein-Trager Resultant

Let $A, B \in \mathbb{K}[x]$ with $\deg(A) < \deg(B)$ and squarefree monic denominator B. The rational function F = A/B has simple poles only.

If
$$F = \sum_{i} \frac{\gamma_i}{x - \beta_i}$$
, then the residue γ_i of F at the pole β_i equals $\gamma_i = \frac{A(\beta_i)}{B'(\beta_i)}$.

Theorem. The residues γ_i of F are roots of the Rothstein-Trager resultant

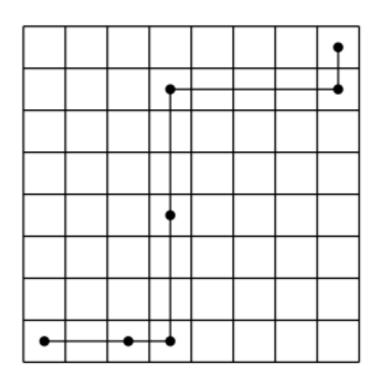
$$R(t) = \operatorname{Res}_{x} (B(x), A(x) - t \cdot B'(x)).$$

Proof. Poisson formula again:
$$R(t) = \prod_{i} (A(\beta_i) - t \cdot B'(\beta_i)).$$

▶ This special resultant is useful for symbolic integration of rational functions.

Application: Diagonal Rook Paths

Question: A chess Rook can move any number of squares horizontally or vertically in one step. How many paths can a Rook take from the lower-left corner square to the upper-right corner square of an $N \times N$ chessboard? Assume that the Rook moves right or up at each step.



Diagonal Rook Paths II

 $1, 2, 14, 106, 838, 6802, 56190, 470010, \dots$

$$\mathsf{Diag}(F) = [s^0] \, F(s, x/s) = \frac{1}{2i\pi} \oint F(s, x/s) \, \frac{ds}{s}, \quad \text{where} \quad F = \frac{1}{1 - \frac{s}{1-s} - \frac{t}{1-t}}.$$

By the residue theorem, Diag(F) is a sum of roots of the Rothstein-Trager resultant

- > F:=1/(1-s/(1-s)-t/(1-t)):
- > G:=normal(1/s*subs(t=x/s,F)):
- > factor(resultant(denom(G),numer(G)-t*diff(denom(G),s),s));

$$x^{2} (2 t - 1) (x - 1) (36 t^{2} x - 4 t^{2} - x + 1)$$

Answer: Generating series of diagonal Rook paths is $\frac{1}{2}\left(1+\sqrt{\frac{1-x}{1-9x}}\right)$.

Application: Certified Algebraic Guessing

Guess+Bound=Proof

Theorem. Suppose $A \in \mathbb{K}[[x]]$ is an algebraic series, and that it is a root of a (unknown) polynomial in $\mathbb{K}[x,y]$ of degree at most d in x and at most n in y.

If
$$\sum_{i=0}^{n} Q_i(x)A^i(x) = O(x^{2dn})$$
, then $\sum_{i=0}^{n} Q_i(x)A^i(x) = 0$.

Proof: Let $P \in \mathbb{K}[x,y]$ be an irreducible polynomial such that

$$P(x, A(x)) = 0$$
, and $\deg_x(P) \le d$, $\deg_y(P) \le n$.

- By Hadamard, $R(x) = \text{Res}_{y}(P,Q) \in \mathbb{K}[x]$ has degree at most 2dn.
- By elimination, R(x) = UP + VQ for $U, V \in \mathbb{K}[x, y]$ with $\deg_y(V) < n$.
- Evaluation at y = A(x) yields

$$R(x) = U(x, A(x)) \underbrace{P(x, A(x))}_{0} + V(x, A(x)) \underbrace{Q(x, A(x))}_{O(x^{2dn})} = O(x^{2dn}).$$

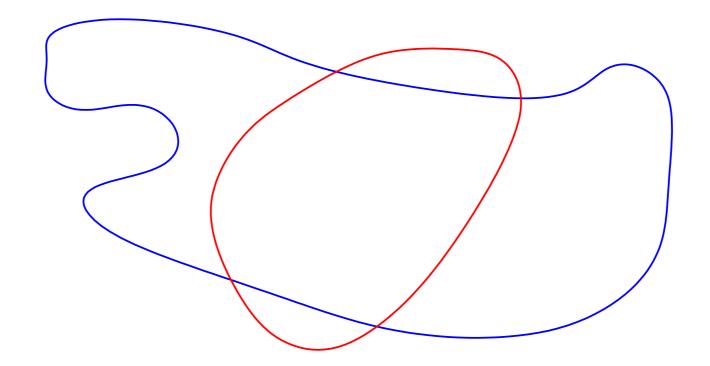
• Thus R = 0, that is $gcd(P, Q) \neq 1$, and thus $P \mid Q$, and A is a root of Q.

Systems of 2 Equations in 2 Unknowns

Geometrically, roots of a polynomial $f \in \mathbb{Q}[x]$ correspond to points on a line.



Roots of polynomials $A \in \mathbb{Q}[x,y]$ correspond to plane curves A=0.

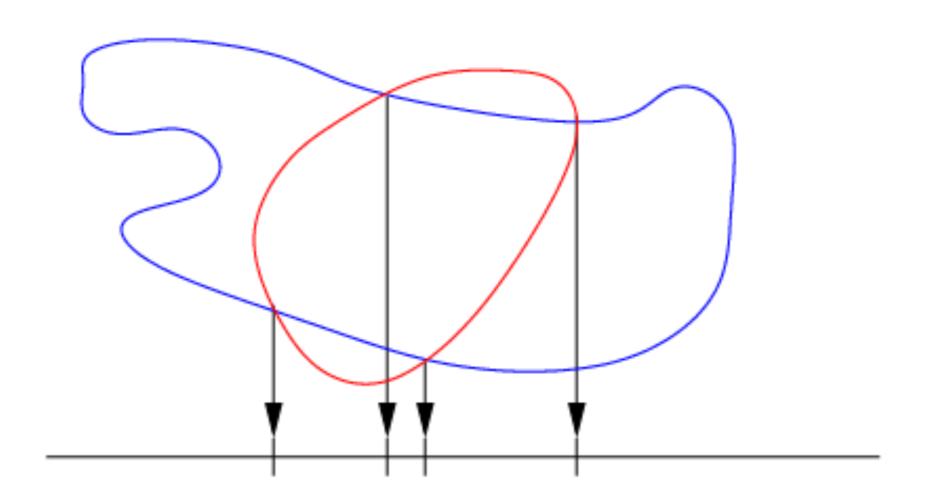


Let now A and B be in $\mathbb{Q}[x,y]$. Then:

- either the curves A = 0 and B = 0 have a common component,
- or they intersect in a finite number of points.

Resultants Compute Projections

Theorem. Let $A = a_m y^m + \cdots$ and $B = b_n y^n + \cdots$ be polynomials in $\mathbb{Q}[x][y]$. The roots of $\text{Res}_y(A, B) \in \mathbb{Q}[x]$ are either the abscissas of points in the intersection A = B = 0, or common roots of a_m and b_n .



Proof. Elimination property: Res (A, B) = UA + VB. Thus $A(\alpha, \beta) = B(\alpha, \beta) = 0$ implies Res $_y(A, B)(\alpha) = 0$

Application: Implicitization of Parametric Curves

Task: Given a rational parametrization of a curve

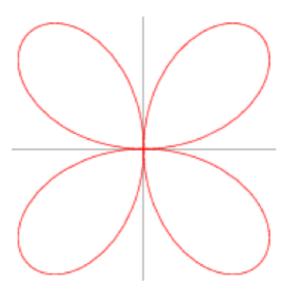
$$x = A(t), \quad y = B(t), \quad A, B \in \mathbb{K}(t),$$

compute a non-trivial polynomial in x and y vanishing on the curve.

Recipe: take the resultant in t of numerators of x - A(t) and y - B(t).

Example: for the four-leaved clover (a.k.a. quadrifolium) given by

$$x = \frac{4t(1-t^2)^2}{(1+t^2)^3}, \quad y = \frac{8t^2(1-t^2)}{(1+t^2)^3},$$



$$\operatorname{Res}_{t}((1+t^{2})^{3}x-4t(1-t^{2})^{2},(1+t^{2})^{3}y-8t^{2}(1-t^{2}))=2^{24}\left((x^{2}+y^{2})^{3}-4x^{2}y^{2}\right).$$

2. D-finite Series and Sequences

Differential or Recurrence equations as a data-structure

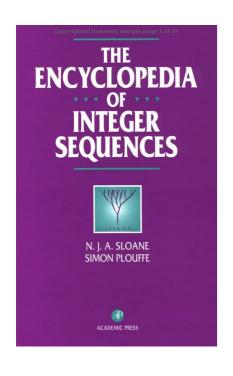
D-finite Series & Sequences

Definition: A power series $f(x) \in \mathbb{K}[[x]]$ is D-finite over \mathbb{K} when its derivatives generate a finite-dimensional vector space over $\mathbb{K}(x)$.

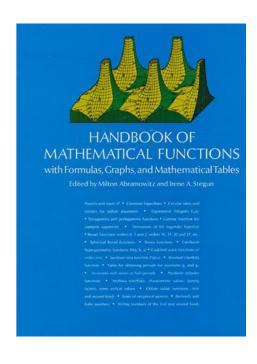
A sequence u_n is D-finite (or P-recursive) over \mathbb{K} when its shifts (u_n, u_{n+1}, \dots) generate a finite-dimensional vector space over $\mathbb{K}(n)$.

equation + init conditions = data structure

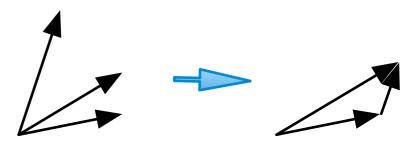
About 25% of Sloane's encyclopedia, 60% of Abramowitz & Stegun



Examples: exp, log, sin, cos, sinh, cosh, arccos, arccosh, arcsin, arcsinh, arctan, arctanh, arccot, arccoth, arccsc, arccsch, arcsec, arcsech, $_pF_q$ (includes Bessel J, Y, I and K, Airy Ai and Bi and polylogarithms), Struve, Weber and Anger functions, the large class of algebraic functions,...



Automatic Proof of Identities



 $> series(sin(x)^2+cos(x)^2-1,x,4);$

f satisfies a LDE ⇔

f,f',f'',... live in a

finite-dim. vector space $O(x^4)$

Why is this a proof?

- 1. sin and cos satisfy a 2nd order LDE: y''+y=0;
- 2. their squares and their sum satisfy a 3rd order LDE;
- 3. the constant -1 satisfies y'=0;
- 4. thus sin²+cos²-1 satisfies a LDE of order at most 4;
- 5. the Cauchy-Lipschitz theorem concludes.

Proofs of non-linear identities by linear algebra!

Mehler's identity for Hermite polynomials

$$\sum_{n=0}^{\infty} H_n(x) H_n(y) \frac{u^n}{n!} = \frac{\exp\left(\frac{4u(xy - u(x^2 + y^2))}{1 - 4u^2}\right)}{\sqrt{1 - 4u^2}}$$

- 1. Definition of Hermite polynomials: recurrence of order 2;
- 2. Product by linear algebra: $H_{n+k}(x)H_{n+k}(y)/(n+k)!$, $k \in \mathbb{N}$ generated over $\mathbb{Q}(x,n)$ by

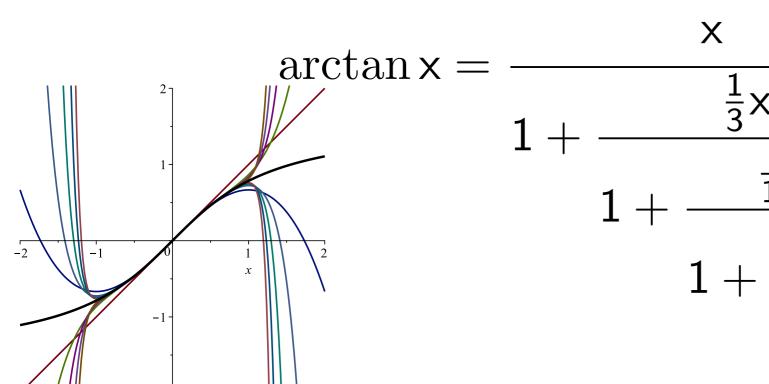
$$\frac{H_n(x)H_n(y)}{n!}, \frac{H_{n+1}(x)H_n(y)}{n!}, \frac{H_n(x)H_{n+1}(y)}{n!}, \frac{H_n(x)H_{n+1}(y)}{n!}, \frac{H_{n+1}(x)H_{n+1}(y)}{n!}$$

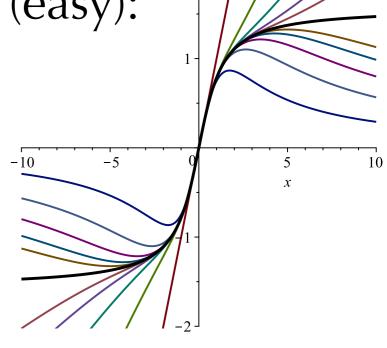
- → recurrence of order at most 4;
- 3. Translate into differential equation.



Guess & prove continued fractions

1. Taylor expansion produces first terms (easy):





- 2. Guess a formula (easy): $a_n = \frac{\pi}{4n^2 1}$
- 3. Prove that the CF with these an converges to arctan.

No human intervention needed.

gfun[ContFrac]

Automatic Proof of the Guessed CF (1/2)

$$\arctan x \stackrel{?}{=} \boxed{1} + \dots + \boxed{\frac{n^2}{4n^2 - 1}x^2} + \dots$$

Lemma. Let P_n/Q_n be the nth convergent. Then

$$\lim_{n \to \infty} \left((x^2 + 1) \left(\frac{P_n}{Q_n} \right)' - 1 \right) = 0 \Rightarrow \lim_{n \to \infty} \frac{P_n}{Q_n} = \arctan x.$$

Let
$$H_n := Q_n^2((x^2 + 1)(P_n/Q_n)' - 1)$$
.

- 1. Compute a linear recurrence for H_n ;
- 2. Using the initial conditions, find a smaller-order one;
- 3. Conclude that $H_n = O(x^n)$.

Automatic Proof of the Guessed CF (2/2)

$$H_n := Q_n^2((x^2+1)(P_n/Q_n)'-1).$$

- 1. P_n and Q_n satisfy $u_n = u_{n-1} + a_n x^2 u_{n-2}$ and $Q_n(0) \neq 0$.
- 2. H_n is a polynomial in P_n , Q_n and their derivatives.
- 3. All its shifts H_{n+k} are linear combinations of $P'_{n+i}Q_{n+j}, P_{n+i}Q'_{n+j}, P_{n+i}P_{n+j}, Q_{n+i}Q_{n+j}, i \text{ and } j \text{ in } \{0,1\}$
- 4. → by linear algebra

$$H_{n+4} = H_{n+3} + (\cdots x^2 + \cdots x^4) H_{n+2} + \cdots x^4 H_{n+1} + \cdots x^8 H_n$$

5. Using the initial conditions gives

$$(2n+3)^2 H_{n+1} + (n+2)^2 x^2 H_n = 0.$$

All these steps are easy to automate.

Algebraic Series can be Computed Fast

$$P(X, Y(X)) = 0$$
 P irreducible

Wanted: the first *N* Taylor coefficients of *Y*.

$$P_x(X,Y(X)) + P_y(X,Y(X)) \cdot Y'(X) = 0$$

$$Y'(X) = (-P_x P_y^{-1} \bmod P)(X,Y(X))$$
a polynomial

Note: F sol LDE $\Rightarrow F(Y(X)) \text{ sol LDE}$ (same argument)

$$Y(X), Y'(X), Y''(X), \dots$$
 in $\operatorname{Vect}_{\mathbb{Q}(X)}(1, Y, Y^2, \dots)$

finite dimension

→ a LDE by linear algebra

An Olympiad Problem

Question: Let (a_n) be the sequence with $a_0 = a_1 = 1$ satisfying the recurrence

$$(n+3)a_{n+1} = (2n+3)a_n + 3na_{n-1}.$$

Show that all a_n is an integer for all n.

Computer-aided solution: Let's compute the first 10 terms of the sequence:

```
> rec:=(n+3)*a(n+1)-(2*n+3)*a(n)-3*n*a(n-1): ini:=a(0)=1,a(1)=1:
> pro:=gfun:-rectoproc({rec,ini}, a(n), list);
```

> pro(10);

gfun's seriestoalgeq command allows to guess that GF is algebraic:

Thus it is very likely that $y = \sum_{n\geq 0} a_n x^n$ verifies $1 + (x-1)y + x^2y^2 = 0$. By coefficient extraction, (a_n) conjecturally verifies the non-linear recurrence

$$a_{n+2} = a_{n+1} + \sum_{k=0}^{n} a_k \cdot a_{n-k}.$$
 (1)

Clearly (1) implies $a_n \in \mathbb{N}$. To prove (1), we proceed the other way around: we start with $P(x,y) = 1 + (x-1)y + x^2y^2$, and show that it admits a power series solution whose coefficients satisfy the same linear recurrence as (a_n) :

- > deq:=gfun:-algeqtodiffeq(pol,y(x)):
- > recb:=gfun:-diffeqtorec(deq,y(x),b(n));

recb :=
$$\{(3 + 3 n) b(n) + (2 n + 5) b(n + 1) + (-4 - n) b(n + 2), b(0) = 1, b(1) = 1\}$$

► In fact,

$$a_n = \sum_{k=0}^n \frac{(-1)^{n-k}}{k+2} \binom{n}{k} \binom{2k+2}{k+1} = \sum_{j\geq 0} (-1)^j \binom{n+1}{j} \binom{2n-3j}{n},$$

(which clearly implies $a_n \in \mathbb{Z}$) but how to find algorithmically such a formula?

3. Creative Telescoping

Examples I: hypergeometric summation

$$\bullet \sum_{k \in \mathbb{Z}} (-1)^k \binom{a+b}{a+k} \binom{a+c}{c+k} \binom{b+c}{b+k} = \frac{(a+b+c)!}{a!b!c!}$$

•
$$A_n = \sum_{k=0}^{n} {n \choose k}^2 {n+k \choose k}^2$$
 satisfies the recurrence [Apéry78]:

$$(n+1)^3 A_{n+1} = (34n^3 + 51n^2 + 27n + 5)A_n - n^3 A_{n-1}.$$

(Neither Cohen nor I had been able to prove this in the intervening two months [Van der Poorten]).

•
$$\sum_{k=0}^{n} {n \choose k}^2 {n+k \choose k}^2 = \sum_{k=0}^{n} {n \choose k} {n+k \choose k} \sum_{j=0}^{k} {n \choose k}^3$$
 [Strehl92]

Examples II: Integrals

•
$$\int_0^1 \frac{\cos(zu)}{\sqrt{1-u^2}} du = \int_1^{+\infty} \frac{\sin(zu)}{\sqrt{u^2-1}} du = \frac{\pi}{2} J_0(z);$$

•
$$\int_0^{+\infty} x J_1(ax) I_1(ax) Y_0(x) K_0(x) dx = -\frac{\ln(1-a^4)}{2\pi a^2}$$
 [Glasser-Montaldi94];

•
$$\frac{1}{2\pi i} \oint \frac{(1+2xy+4y^2) \exp\left(\frac{4x^2y^2}{1+4y^2}\right)}{y^{n+1}(1+4y^2)^{\frac{3}{2}}} dy = \frac{H_n(x)}{\lfloor n/2 \rfloor!}$$
 [Doetsch30].

Examples III: Diagonals

Definition If
$$f(x_1, \ldots, x_k) = \sum_{i_1, i_2, \ldots, i_k \geq 0} c_{i_1, \ldots, i_k} x_1^{i_1} \cdots x_k^{i_k} \in \mathbb{K}[[x_1, \ldots, x_k]]$$
, then its diagonal is $\operatorname{Diag}(f) = \sum_{n \geq 0} c_{n, \ldots, n} x^n \in \mathbb{K}[[x]]$.

- Diagonal k-D rook paths: Diag $\frac{1}{1 \frac{x_1}{1 x_1} \dots \frac{x_k}{1 x_k}}$;
- Hadamard product: $F(x) \odot G(x) = \sum_n f_n g_n x^n = \text{Diag}(F(x)G(y));$
- Algebraic series [Furstenberg67]: if P(x, S(x)) = 0 and $P_y(0, 0) \neq 0$ then

$$S(x) = \text{Diag}\left(y^2 \frac{P_y(xy,y)}{P(xy,y)}\right).$$

• Apéry's sequence [Dwork80]:

$$\sum A_n z^n = \text{Diag} \frac{1}{(1-x_1)((1-x_2)(1-x_3)(1-x_4)(1-x_5)-x_1x_2x_3)}.$$

Theorem [Lipshitz88] The diagonal of a rational (or algebraic, or even D-finite) series is D-finite.

Summation by Creative Telescoping

$$I_n := \sum_{k=0}^n \binom{n}{k} = 2^n.$$

IF one knows Pascal's triangle:

$$\binom{n+1}{k} = \binom{n}{k} + \binom{n}{k-1} = 2\binom{n}{k} + \binom{n}{k-1} - \binom{n}{k},$$

then summing over k gives

$$I_{n+1}=2I_n.$$

The initial condition $I_0 = 1$ concludes the proof.

Creative Telescoping for Sums

$$F_n = \sum_k u_{n,k} = ?$$

IF one knows $A(n, S_n)$ and $B(n, k, S_n, S_k)$ s.t.

$$(A(n, S_n) + \Delta_k B(n, k, S_n, S_k)) \cdot u_{n,k} = 0$$

(where Δ_k is the difference operator, $\Delta_k \cdot v_{n,k} = v_{n,k+1} - v_{n,k}$), then the sum "telescopes", leading to

$$A(n, S_n) \cdot F_n = 0.$$

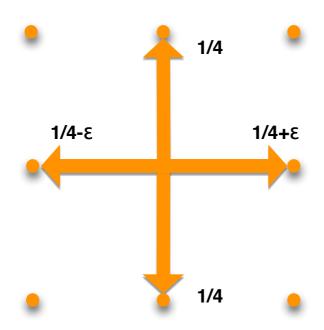
Zeilberger's Algorithm [1990]

Input: a hypergeometric term $u_{n,k}$, i.e., $u_{n+1,k}/u_{n,k}$ and $u_{n,k+1}/u_{n,k}$ rational functions in n and k;

Output:

- a linear recurrence (A) satisfied by $F_n = \sum_k u_{n,k}$
- a certificate (B), s.t. checking the result is easy from $A(n, S_n) \cdot u_{n,k} = \Delta_k B \cdot u_{n,k}$.

Example: SIAM flea



$$U_{n,k} := {2n \choose 2k} {2k \choose k} {2n-2k \choose n-k} \left(\frac{1}{4}+c\right)^k \left(\frac{1}{4}-c\right)^k \frac{1}{4^{2n-2k}}.$$

> SumTools[Hypergeometric][Zeilberger](U,n,k,Sn);

$$[(4n^{2} + 16n + 16) Sn^{2} + (-4n^{2} + 32c^{2}n^{2} + 96c^{2}n - 12n + 72c^{2} - 9) Sn + 128c^{4}n + 64c^{4}n^{2} + 48c^{4}, ... (BIG certificate)...]$$

Creative Telescoping for Integrals

$$I(x) = \int_{\Omega} u(x, y) \, dy = ?$$

IF one knows $A(x, \partial_x)$ and $B(x, y, \partial_x, \partial_y)$ s.t.

$$(A(x, \partial_x) + \partial_y B(x, y, \partial_x, \partial_y)) \cdot u(x, y) = 0,$$

then the integral "telescopes", leading to

$$A(x, \partial_x) \cdot I(x) = 0.$$

Special Case: Diagonals

Analytically,

$$Diag(F(x,y)) = \frac{1}{2\pi i} \oint F\left(\frac{x}{y}, y\right) \frac{dy}{y}.$$

On power series,

$$(A(x, \partial_x) + \partial_y B) \cdot \underbrace{\frac{1}{y} F\left(\frac{x}{y}, y\right)}_{U} = 0 \Longrightarrow A(x, \partial_x) \cdot \text{Diag } F = 0.$$

Proof:

1.
$$[y^{-1}]U = \text{Diag}(f);$$

2.
$$[y^{-1}]A \cdot U + [y^{-1}]\partial_y B \cdot U = A \cdot [y^{-1}]U$$
.

Extends to more variables: Diag F(x, y, z) obtained from $[y^{-1}z^{-1}]U$, $U = \frac{1}{yz}F\left(\frac{x}{y}, \frac{y}{z}, z\right)$, if one finds

$$(A(x, \partial_x) + \partial_y B(x, y, z, \partial_x, \partial_y, \partial_z) + \partial_z C(x, y, z, \partial_x, \partial_y, \partial_z)) \cdot U = 0.$$

Provided by Chyzak's algorithm

Summary of the Exercises for this Afternoon

4. A nice number

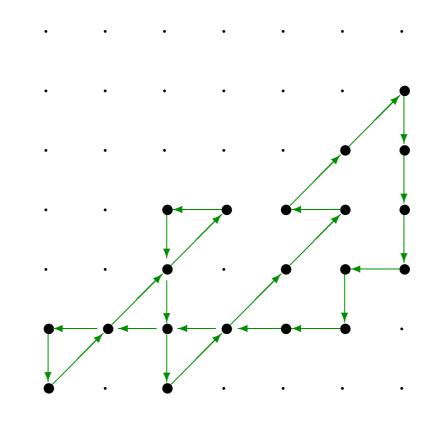
Guess and prove a simple formula for $\frac{\sin\frac{2\pi}{7}}{\sin^2\frac{3\pi}{7}} - \frac{\sin\frac{\pi}{7}}{\sin^2\frac{2\pi}{7}} + \frac{\sin\frac{3\pi}{7}}{\sin^2\frac{\pi}{7}}$. See also the exercises on resultants.

$$\frac{\sin\frac{2\pi}{7}}{\sin^2\frac{3\pi}{7}} - \frac{\sin\frac{\pi}{7}}{\sin^2\frac{2\pi}{7}} + \frac{\sin\frac{3\pi}{7}}{\sin^2\frac{\pi}{7}}$$

Pb. Kreweras Excursions

$$K(t; x, y) = \sum_{n=0}^{\infty} \left(\sum_{i,j \ge 0} k(n; i, j) x^i y^j \right) t^n$$

is algebraic.



THE END

(Except for the exercises!)